

四值 Gödel 命题逻辑系统中公式的概率真度理论

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摘要 在四值 Gödel 命题逻辑系统中提出了公式的概率真度,证明了全体公式的概率真度值之集在 $[0,1]$ 中没有孤立点;定义了两个公式间的概率相似度,建立了概率逻辑度量空间,证明了此空间中没有孤立点,为研究四值 Gödel 命题逻辑系统的近似推理提供了思路。

关键词 概率真度,相似度,孤立点

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Theory of Probability Truth Degree in Gödel 4-valued Propositional Logic System

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Abstract This paper proposed the theory of probability truth degree in Gödel 4-valued propositional logic. It is proved that the set of probability truth degree of all formulas has no isolated point in $[0,1]$. The conceptions of probability similarity degree on two formulas were defined. Moreover, the probability logic metric space was built. It was proved that this space has no isolated point. Then it can provide the thought for approximate reasoning theory.

Keywords Probability truth degree, Similarity degree, Isolated point

1 引言

在等概率即每个原子公式 p 赋值为 0 和 1 的概率均为 $\frac{1}{2}$ 的情况下,王国俊等^[1-6]研究了均匀概率空间的情形;文献[7]中,李骏等将公式真度概念推广到三值 Łukasiewicz 逻辑系统中,研究了每个原子公式 p 赋值为 $0, \frac{1}{2}, 1$ 的概率测度均为 $\frac{1}{3}$ 的情形,但这种等概率的研究不能很好地体现概率具有随机性的特点。

在非均匀概率空间中,关晓红等^[8,9]利用概率测度的方法,研究了三值 Łukasiewicz 逻辑系统以及标准序列逻辑系统 S_3 中原子公式 $q_i (i=1,2,\dots)$ 随机赋值为 $0, \frac{1}{2}, 1$ 的概率测度分别为 $p_{i0}, p_{i\frac{1}{2}}, p_{i1}$ 且 $p_{i0} + p_{i\frac{1}{2}} + p_{i1} = 1$ 的情形。本文利用文献[8,9]的思想,在四值 Gödel 命题逻辑系统中提出了公式的概率真度,证明了所有公式的概率真度值之集在 $[0,1]$ 中没有孤立点;定义了两个公式间的概率相似度,建立了概率逻辑度量空间,证明了此空间中没有孤立点,因此为研究四值 Gödel 命题逻辑系统近似推理提供了思路。

2 逻辑公式的概率真度

定义 1^[10] 设 (X_n, A_n, μ_n) 是概率测度空间, μ_n 是 X_n 上的 Lebesgue 测度, A_n 是全体 μ_n 可测集之族且 $\mu_n(X_n) = 1$ 。

令 $X = \prod_{n=1}^{\infty} X_n$, 则 $\prod_{n=1}^{\infty} X_n$ 在 X 上可生成一个 σ -代数 A 。这时 X 上存在唯一的测度 μ 满足以下条件:

(1) A 是 X 中的可测集之族;

(2) 对于 $\prod_{n=1}^{\infty} X_n$ 中的任一可测集 $E, E \times \prod_{n=m+1}^{\infty} X_n$ 可测,且

$$\mu(E \times \prod_{n=m+1}^{\infty} X_n) = (\mu_1 \times \mu_2 \times \dots \times \mu_m)(E), m=1,2,\dots$$

称 μ 为 X 上的关于 μ_1, μ_2, \dots 的无穷乘积测度。

定义 2^[10] 设 n 是固定的自然数, $n \geq 2, (Y, B, \eta)$ 是非均匀概率测度空间, 这里 $Y = \{y_1, y_2, \dots, y_n\}, \eta(\Phi) = 0, \eta(Y) = 1$ 。令 $(X_k, A_k, \mu_k) = (Y, B, \mu) (k=1,2,\dots)$, 且 (X, A, μ) 是 $\{(X_k, A_k, \mu_k)\}_{k=1}^{\infty}$ 的无穷乘积, 称 (X, A, μ) 是 n 值逻辑概率测度空间。逻辑概率测度空间 (X, A, μ) 也可记为 X 。

设 $S = \{q_1, q_2, \dots\}$ (为了与概率值符号 p 区分, 本文原子公式用 q 来表示) 为原子公式集, N 是由 S 生成的 $(\neg, \vee, \rightarrow)$ 型自由代数, 这里 \neg, \vee, \rightarrow 分别是一元和二元逻辑连接词, 称 S 中的元素为原子命题(或原子公式), 称 $F(S)$ 中的元素为命题(或公式)。

令 $L_4 = \{0, \frac{1}{3}, \frac{2}{3}, 1\}$, 则 $X_n = L_4 = \{0, \frac{1}{3}, \frac{2}{3}, 1\}$ 。在 L_4 中规定 $\forall x, y \in L_4, \neg x = x \rightarrow 0, x \vee y = \max\{x, y\}, x \rightarrow y =$

$$R_G(x, y) (R_G(x, y) = \begin{cases} 1, & x \leq y \\ 0, & x > y \end{cases}), 则 L 成为 $(\neg, \vee, \rightarrow)$ 型的代数。$$

$\forall A \in F(S), v(A)$ 叫做公式 A 的赋值。 $F(S)$ 的赋值映射的全体记为 Ω 。

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对任一公式 $A \in F(S)$, 以下把公式 $(\neg A \rightarrow A)$ 简记为 $2A$.
容易验证

$$\forall v \in \Omega, v(2A) = \begin{cases} 0, & v(A) = 0 \\ 1, & v(A) = \frac{1}{3}, \frac{2}{3}, 1 \end{cases}$$

定义 3 设 $v \in \Omega$, 则根据 $F(S)$ 是由生成的自由代数知 v 由 $v|S$ 唯一确定. 设 $v(q_k) = v_k (k=1, 2, \dots)$, 则无穷维向量 $v = (v_1, v_2, \dots) \in X$, 这里 X 由定义 2 确定. 反过来, 设 $v = (v_1, v_2, \dots) \in X$, 则 v 唯一确定 Ω 中的一个赋值 v , 这里 $v(q_k) = v_k (k=1, 2, \dots)$. 令 $\varphi(v) = \nu$, 则 $\varphi: \Omega \rightarrow X$ 是从 Ω 到 X 的一一满射, 称 φ 为 Ω 的测度化映射.

定义 4 设 $X_n = \{0, \frac{1}{3}, \frac{2}{3}, 1\}$, μ_n 是 X_n 上的非均匀的概率测度, 即 $\mu_n(\Phi) = 0, \mu_n(Y) = 1, \mu_n\{v \in X | v \in \Omega, v(q_n) = \frac{i}{3}\} = p_{n,i} (0 < p_{n,i} < 1$ 且 $p_{n,0} + p_{n,\frac{1}{3}} + p_{n,\frac{2}{3}} + p_{n,1} = 1; i=0, 1, 2, 3; n=1, 2, \dots)$, 令 $X = \prod_{n=1}^{\infty} X_n$, 设 μ 是 X 上的关于 μ_1, μ_2, \dots 的无穷测度, 称 μ 为四值逻辑概率测度.

定义 5 设 $A \in F(S)$, 令 $[A_{\frac{i}{3}}] = \{v \in X | v \in \Omega, v(A) = \frac{i}{3}\} (i=0, 1, 2, 3)$,

$$\tau^*(A) = \sum_{i=0}^3 \frac{i}{3} \mu([A_{\frac{i}{3}}])$$

称 $\tau^*(A)$ 为逻辑公式 A 的概率真度.

显然, 逻辑等价的公式具有相同的概率真度.

例 1 设 $A = q_1, B = q_1 \rightarrow q_2, C = q_1 \wedge q_2$, 计算 $\tau^*(A), \tau^*(B), \tau^*(C)$. 这里假设 $p_{10} = 0.2, p_{1,\frac{1}{3}} = 0.3, p_{1,\frac{2}{3}} = 0.1, p_{11} = 0.4, p_{20} = 0.4, p_{2,\frac{1}{3}} = 0.2, p_{2,\frac{2}{3}} = 0.2, p_{21} = 0.2$.

解: (1) 由 $\mu([A_1]) = \mu(\{v \in X | v \in \Omega, v(q_1) = 1\}) = 0.4$;
 $\mu([A_{\frac{2}{3}}]) = \mu(\{v \in X | v \in \Omega, v(q_1) = \frac{2}{3}\}) = 0.1$; $\mu([A_{\frac{1}{3}}]) = \mu(\{v \in X | v \in \Omega, v(q_1) = \frac{1}{3}\}) = 0.3$ 得: $\tau^*(A) = \frac{1}{3} \mu([A_{\frac{1}{3}}]) + \frac{2}{3} \mu([A_{\frac{2}{3}}]) + \mu([A_1]) = \frac{1}{3} \times 0.3 + \frac{2}{3} \times 0.1 + 0.4 = \frac{17}{30}$.

(2) 由 $[B_1] = \{v \in X | v \in \Omega, v(q_1 \rightarrow q_2) = 1\} = \{v \in X | v(q_2) = 1\} \cup \{v \in X | v(q_1) = v(q_2) = \frac{2}{3}\} \cup \{v \in X | v(q_1) = v(q_2) = \frac{1}{3}\} \cup \{v \in X | v(q_1) = v(q_2) = 0\} \cup \{v \in X | v(q_1) = 0, v(q_2) = \frac{1}{3}\} \cup \{v \in X | v(q_1) = 0, v(q_2) = \frac{2}{3}\} \cup \{v \in X | v(q_1) = \frac{1}{3}, v(q_2) = \frac{2}{3}\}$; $[B_{\frac{2}{3}}] = \{v \in X | v \in \Omega, v(q_1 \rightarrow q_2) = \frac{2}{3}\} = \{v \in X | v(q_1) = 1, v(q_2) = \frac{2}{3}\}$; $[B_{\frac{1}{3}}] = \{v \in X | v \in \Omega, v(q_1 \rightarrow q_2) = \frac{1}{3}\} = \{v \in X | v(q_1) = 1, v(q_2) = \frac{1}{3}\} \cup \{v \in X | v(q_1) = \frac{2}{3}, v(q_2) = \frac{1}{3}\}$ 得: $\tau^*(B) = \frac{1}{3} \mu([B_{\frac{1}{3}}]) + \frac{2}{3} \mu([B_{\frac{2}{3}}]) + \mu([B_1]) = \frac{1}{3} \times (0.4 \times 0.2 + 0.1 \times 0.2) + \frac{2}{3} \times 0.4 \times 0.2 + 0.2 + 0.1 \times 0.2 + 0.3 \times 0.2 + 0.2 \times 0.4 + 0.2 \times 0.2 + 0.2 \times 0.2 + 0.3 \times 0.2 = \frac{44}{75}$.

(3) 由 $[C_1] = \{v \in X | v \in \Omega, v(q_1) = v(q_2) = 1\}$; $[C_{\frac{2}{3}}] = \{v \in X | v \in \Omega, v(q_1) = 1, v(q_2) = \frac{2}{3}\} \cup \{v \in X | v \in \Omega, v(q_1) = \frac{2}{3}, v(q_2) = 1\}$; $[C_{\frac{1}{3}}] = \{v \in X | v \in \Omega, v(q_1) = 1, v(q_2) = \frac{1}{3}\} \cup \{v \in X | v \in \Omega, v(q_1) = \frac{1}{3}, v(q_2) = 1\} \cup \{v \in X | v \in \Omega, v(q_1) = \frac{2}{3}, v(q_2) = \frac{1}{3}\} \cup \{v \in X | v \in \Omega, v(q_1) = \frac{1}{3}, v(q_2) = \frac{2}{3}\}$ 得: $\tau^*(C) = \frac{1}{3} \mu([C_{\frac{1}{3}}]) + \frac{2}{3} \mu([C_{\frac{2}{3}}]) + \mu([C_1]) = \frac{1}{3} \times (0.4 \times 0.2 + 0.2 + 0.3 \times 0.2 + 0.1 \times 0.2 + 0.3 \times 0.2) + \frac{2}{3} \times (0.4 \times 0.2 + 0.1 \times 0.2) + 0.4 \times 0.2 = \frac{11}{50}$.

定理 1 设 $A \in F(S)$, 则

(1) A 是重言式当且仅当 $\tau^*(A) = 1$;

(2) A 是矛盾式当且仅当 $\tau^*(A) = 0$.

证明: 设 $A = A(q_1, q_2, \dots, q_n)$ 是重言式, 则 $[A_1] = \{v \in X | v(A) = 1\}, [A_{\frac{2}{3}}] = [A_{\frac{1}{3}}] = [A_0] = \Phi$, 从而 $\tau^*(A) = 1$. 反之, 设 A 不是重言式, 则有 $v \in \Omega$, 使 $v(A) = 0, \frac{1}{3}, \frac{2}{3}$, 即 $\tau^*(A) = 0 \times \mu([A_0]) + \frac{1}{3} \mu([A_{\frac{1}{3}}]) + \frac{2}{3} \mu([A_{\frac{2}{3}}]) + \mu([A_1]) < \mu([A_0]) + \mu([A_{\frac{1}{3}}]) + \mu([A_{\frac{2}{3}}]) + \mu([A_1]) = \mu(X) = 1$. 矛盾式可用以上方法类似得证.

定理 2 设 A 是重言式, 则 $\tau^*(A \rightarrow B) = \tau^*(B), \tau^*(B \rightarrow A) = 1$.

证明: 设 A 是重言式, 则 $v(A) = 1$, 故 $v(A) \rightarrow v(B) = 1 \rightarrow v(B) = \frac{1}{3}$ 当且仅当 $v(B) = \frac{1}{3}$; $v(A) \rightarrow v(B) = 1 \rightarrow v(B) = \frac{2}{3}$ 当且仅当 $v(B) = \frac{2}{3}$; $v(A) \rightarrow v(B) = 1 \rightarrow v(B) = 1$ 当且仅当 $v(B) = 1$; 而 B 的值不论怎样取, $v(B \rightarrow A) = 1$, 故 $\tau^*(A \rightarrow B) = \frac{1}{3} \mu([(A \rightarrow B)_{\frac{1}{3}}]) + \frac{2}{3} \mu([(A \rightarrow B)_{\frac{2}{3}}]) + \mu([(A \rightarrow B)_1]) = \frac{1}{3} \mu([B_{\frac{1}{3}}]) + \frac{2}{3} \mu([B_{\frac{2}{3}}]) + \mu([B_1]) = \tau^*(B)$; $\tau^*(B \rightarrow A) = \frac{1}{3} \mu([(B \rightarrow A)_{\frac{1}{3}}]) + \frac{2}{3} \mu([(B \rightarrow A)_{\frac{2}{3}}]) + \mu([(B \rightarrow A)_1]) = \frac{1}{3} \mu(\Phi) + \frac{2}{3} \mu(\Phi) + \mu(X) = 1$.

故结论成立.

定理 3 设 $A, B \in F(S)$, 则

(1) $\tau^*(A \vee B) = \tau^*(A) + \tau^*(B) - \tau^*(A \wedge B)$;

(2) 如果 $A = A(q_1, q_2, \dots, q_n), B = B(q_{n+1}, q_{n+2}, \dots, q_{n+m})$, 则 $\tau^*(A) \times \tau^*(B) \leq \tau^*(A \wedge B)$.

证明: (1) $[(A \vee B)_1] = \{v \in X | v(A) = 1 \text{ 或 } v(B) = 1\} = \{v \in X | v(A) = 1\} \cup \{v \in X | v(B) = 1\} = [A_1] \cup [B_1]$; $[(A \vee B)_{\frac{2}{3}}] = \{v \in X | v(A) = v(B) = \frac{2}{3}\} \cup \{v \in X | v(A) = 0, v(B) = \frac{2}{3}\} \cup \{v \in X | v(A) = \frac{2}{3}, v(B) = 0\} \cup \{v \in X | v(A) = \frac{1}{3}, v(B) = \frac{2}{3}\} \cup \{v \in X | v(A) = \frac{2}{3}, v(B) = \frac{1}{3}\}$; $[(A \vee B)_{\frac{1}{3}}] = \{v \in X |$

$$v(A) = v(B) = \frac{1}{3} \} \cup \{ \nu \in X \mid v(A) = 0, v(B) = \frac{1}{3} \} \cup \{ \nu \in X \mid v(A) = \frac{1}{3}, v(B) = 0 \}.$$

$$\begin{aligned} [(A \wedge B)_1] &= \{ \nu \in X \mid v(A) = 1 \text{ 且 } v(B) = 1 \} = \{ \nu \in X \mid v(A) = 1 \} \cap \{ \nu \in X \mid v(B) = 1 \} = [A_1] \cap [B_1]; [(A \wedge B)_{\frac{2}{3}}] = \\ &= \{ \nu \in X \mid v(A) = v(B) = \frac{2}{3} \} \cup \{ \nu \in X \mid v(A) = 1, v(B) = \frac{2}{3} \} \cup \\ &= \{ \nu \in X \mid v(A) = \frac{2}{3}, v(B) = 1 \}; [(A \wedge B)_{\frac{1}{3}}] = \{ \nu \in X \mid v(A) = \\ &= v(B) = \frac{1}{3} \} \cup \{ \nu \in X \mid v(A) = 1, v(B) = \frac{1}{3} \} \cup \{ \nu \in X \mid v(A) = \\ &= \frac{1}{3}, v(B) = 1 \} \cup \{ \nu \in X \mid v(A) = \frac{1}{3}, v(B) = \frac{2}{3} \} \cup \{ \nu \in X \mid v(A) = \\ &= \frac{2}{3}, v(B) = \frac{1}{3} \}. \end{aligned}$$

$$\begin{aligned} \tau^*(A) + \tau^*(B) - \tau^*(A \wedge B) &= \frac{1}{3} \mu([A_{\frac{1}{3}}]) + \frac{2}{3} \mu \\ &([A_{\frac{2}{3}}]) + \mu([A_1]) + \frac{1}{3} \mu([B_{\frac{1}{3}}]) + \frac{2}{3} \mu([B_{\frac{2}{3}}]) + \mu([B_1]) - \\ &= \frac{1}{3} \mu([(A \wedge B)_{\frac{1}{3}}]) - \frac{2}{3} \mu([(A \wedge B)_{\frac{2}{3}}]) - \mu([(A \wedge B)_1]). \end{aligned}$$

由测度的可加性可得:

$$\begin{aligned} \mu([(A \vee B)_1]) &= \mu([A_1]) + \mu([B_1]) - \mu([(A \wedge B)_1]); \\ \mu([A_{\frac{2}{3}}]) + \mu([B_{\frac{2}{3}}]) - \mu([(A \wedge B)_{\frac{2}{3}}]) &= \mu([A_{\frac{2}{3}}]) + \\ &= \mu([B_{\frac{2}{3}}]) - \mu(\{ \nu \in X \mid v(A) = v(B) = \frac{2}{3} \}) - \mu(\{ \nu \in X \mid v(A) = \\ &= 1, v(B) = \frac{2}{3} \}) - \mu(\{ \nu \in X \mid v(A) = \frac{2}{3}, v(B) = 1 \}) = \mu(\{ \nu \in X \mid \\ &= v(A) = v(B) = \frac{2}{3} \}) + \mu(\{ \nu \in X \mid v(A) = \frac{2}{3}, v(B) = \frac{1}{3} \}) + \\ &= \mu(\{ \nu \in X \mid v(A) = \frac{2}{3}, v(B) = 0 \}) + \mu(\{ \nu \in X \mid v(A) = \frac{1}{3}, v(B) = \\ &= \frac{2}{3} \}) + \mu(\{ \nu \in X \mid v(A) = 0, v(B) = \frac{2}{3} \}). \end{aligned}$$

$$\begin{aligned} \text{同理 } \mu([A_{\frac{1}{3}}]) + \mu([B_{\frac{1}{3}}]) - \mu([(A \wedge B)_{\frac{1}{3}}]) &= \mu(\{ \nu \in X \mid \\ &= v(A) = v(B) = \frac{1}{3} \}) + \mu(\{ \nu \in X \mid v(A) = \frac{1}{3}, v(B) = 0 \}) + \mu(\{ \nu \in \\ &= X \mid v(A) = 0, v(B) = \frac{1}{3} \}). \end{aligned}$$

比较以上各项可得: $\tau^*(A \vee B) = \tau^*(A) + \tau^*(B) - \tau^*(A \wedge B)$.

$$\begin{aligned} (2) [(A \wedge B)_{\frac{1}{3}}] &= \{ \nu \in X \mid v(A) = v(B) = \frac{1}{3} \text{ 或 } v(A) = \\ &= \frac{1}{3}, v(B) = 1 \text{ 或 } v(A) = 1, v(B) = \frac{1}{3} \text{ 或 } v(A) = \frac{2}{3}, v(B) = \frac{1}{3} \\ &= \text{ 或 } v(A) = \frac{1}{3}, v(B) = \frac{2}{3} \}; [(A \wedge B)_{\frac{2}{3}}] = \{ \nu \in X \mid v(A) = v(B) = \\ &= \frac{2}{3} \text{ 或 } v(A) = \frac{2}{3}, v(B) = 1 \text{ 或 } v(A) = 1, v(B) = \frac{2}{3} \}; [(A \wedge \\ &= B)_1] = \{ \nu \in X \mid v(A) = v(B) = 1 \}. \end{aligned}$$

$$\begin{aligned} \tau^*(A \wedge B) &= \frac{1}{3} \mu([(A \wedge B)_{\frac{1}{3}}]) + \frac{2}{3} \mu([(A \wedge B)_{\frac{2}{3}}]) + \\ &= \mu([(A \wedge B)_1]) = \mu([A_1])\mu([B_1]) + \frac{1}{3} (\mu([A_{\frac{1}{3}}])\mu([B_{\frac{1}{3}}]) + \\ &= \mu([A_1])\mu([B_{\frac{1}{3}}]) + \mu([A_{\frac{1}{3}}])\mu([B_1]) + \mu([A_{\frac{2}{3}}])\mu([B_{\frac{1}{3}}]) + \\ &= \mu([A_{\frac{1}{3}}])\mu([B_{\frac{2}{3}}])) + \frac{2}{3} (\mu([A_{\frac{2}{3}}])\mu([B_{\frac{2}{3}}]) + \mu([A_{\frac{2}{3}}]) \\ &= \mu([B_1]) + \mu([A_1])\mu([B_{\frac{2}{3}}])) + \mu([A_1])\mu([B_1]). \end{aligned}$$

$$\begin{aligned} \tau^*(A) \times \tau^*(B) &= (\frac{1}{3} \mu([A_{\frac{1}{3}}]) + \frac{2}{3} \mu([A_{\frac{2}{3}}]) + \\ &= \mu([A_1])) (\frac{1}{3} \mu([B_{\frac{1}{3}}]) + \frac{2}{3} \mu([B_{\frac{2}{3}}]) + \mu([B_1])) = \frac{1}{9} \mu \\ &= [A_{\frac{1}{3}}] \mu([B_{\frac{1}{3}}]) + \frac{4}{9} \mu([A_{\frac{2}{3}}]) \mu([B_{\frac{2}{3}}]) + \frac{2}{9} (\mu([A_{\frac{1}{3}}]) \mu([B_{\frac{2}{3}}]) + \mu([A_{\frac{2}{3}}]) \\ &= \mu([B_{\frac{1}{3}}]) + \frac{1}{3} (\mu([A_{\frac{1}{3}}]) \mu([B_1]) + \mu([A_1]) \mu([B_{\frac{1}{3}}]) + \frac{2}{3} (\mu([A_{\frac{2}{3}}]) [B_1] + \\ &= \mu([A_1]) [B_{\frac{2}{3}}]) + \mu([A_1]) [B_1] \leq \tau^*(A \wedge B), \text{ 故(2)成立.} \end{aligned}$$

以下研究概率真度在 $[0, 1]$ 中的分布情形, 首先介绍两个引理:

引理 1^[11] 设有数列 $\{a_n\}$, $\forall n \in \mathbb{N}_+$, $a_n \in (0, 1)$, ϵ 是任给的正数, 如果对任意的 $n_0 \in \mathbb{N}_+$, 均有 $a_1 \times a_2 \times \cdots \times a_{n_0} \geq \epsilon$, 则数列 $\{a_n\}$ 有一个递增的子数列.

引理 2^[11] 设有数列 $\{a_n\}$, $\forall n \in \mathbb{N}_+$, $a_n \in (0, 1)$, 且 $\{a_n\}$ 为递减数列, $0 < \epsilon < 1$, 则存在 $n_0 \in \mathbb{N}_+$, 使得 $a_1 \times a_2 \times \cdots \times a_{n_0} < \epsilon$.

定理 4 $F(S)$ 中的公式的概率真度值之集 $\{\tau^*(A) \mid A \in F(S)\}$ 在 $[0, 1]$ 中没有孤立点.

证明: 设 $A \in F(S)$, $\epsilon > 0$, 下面证明存在 $B \in F(S)$, 使得 $|\tau^*(A) - \tau^*(B)| < \epsilon$, 且 $\tau^*(B) \neq \tau^*(A)$.

(1) 设 $\tau^*(A) = 1$, 令 $B = B(n) = 2q_1 \vee 2q_2 \vee \cdots \vee 2q_n$, 因为只有当 $v(q_1) = v(q_2) = \cdots = v(q_n) = 0$ 时, $v(2q_1 \vee 2q_2 \vee \cdots \vee 2q_n) = 0$, 否则 $v(2q_1 \vee 2q_2 \vee \cdots \vee 2q_n) = 1$, 所以 $\tau^*(\bar{B}(n)) = \tau^*(2q_1 \vee \cdots \vee 2q_n) = 1 - p_{10} \times \cdots \times p_{n0}$, $|\tau^*(A) - \tau^*(B)| = 1 - \tau^*(\bar{B}(n)) = p_{10} \times \cdots \times p_{n0}$.

① 假设存在 n , 使得 $p_{10} \times \cdots \times p_{n0} < \epsilon$, 则 $|\tau^*(A) - \tau^*(B)| < \epsilon$, 由于 $\tau^*(B(n)) = 1 \times (p_{10} \times \cdots \times p_{n0}) \neq 1$, 因此 $B = B(n) = 2q_1 \vee \cdots \vee 2q_n$ 满足.

② 假设对 $\forall n \in \mathbb{N}_+$, 都有 $p_{10} \times \cdots \times p_{n0} \geq \epsilon$, 根据引理 1, 数列 $\{p_{n0}\}$ 有一个递增的子数列 $\{p_{k_0}\}$ ($i \in \mathbb{N}_+$, $k \in \mathbb{N}_+$), 则 $\{1 - p_{k_0}\}$ 为递减数列. 根据引理 2, 存在 $n_0 \in \mathbb{N}_+$, 使得数列 $\{1 - p_{k_0}\}$ 的前 n_0 项之积小于 ϵ , 这时取 $B = \bar{B}(n_0) = (\neg 2q_1) \vee \cdots \vee (\neg 2q_{i_{n_0}})$, 只有 $v(q_1), \dots, v(q_{n_0})$ 全不取 0 时 $v(\bar{B}(n_0)) = 0$, 否则 $v(\bar{B}(n_0)) = 1$. 所以 $\tau^*(\bar{B}(n_0)) = 1 - 1 \times (1 - p_{1_0}) \times \cdots \times (1 - p_{i_{n_0}}) < 1$, $|\tau^*(A) - \tau^*(B)| = 1 - \tau^*(\bar{B}(n_0)) = 1 \times (1 - p_{1_0}) \times \cdots \times (1 - p_{i_{n_0}}) < \epsilon$, 因此 $B = \bar{B}(n_0) = (\neg 2q_1) \vee \cdots \vee (\neg 2q_{i_{n_0}})$ 满足.

(2) 设 $\tau^*(A) = 0$, 令 $B = B(n) = (\neg 2q_1) \wedge (\neg 2q_2) \wedge \cdots \wedge (\neg 2q_n)$, $|\tau^*(A) - \tau^*(B)| = \tau^*(B(n)) = \tau^*((\neg 2q_1) \wedge (\neg 2q_2) \wedge \cdots \wedge (\neg 2q_n))$.

当 $v(q_1) = \cdots = v(q_n) = 0$ 时, $v((\neg 2q_1) \wedge (\neg 2q_2) \wedge \cdots \wedge (\neg 2q_n)) = 1$, 否则 $v((\neg 2q_1) \wedge (\neg 2q_2) \wedge \cdots \wedge (\neg 2q_n)) = 0$, 所以 $|\tau^*(A) - \tau^*(B)| = 1 \times (p_{10} \times \cdots \times p_{n0})$.

① 假设存在 n , 使得 $p_{10} \times \cdots \times p_{n0} < \epsilon$, 则 $|\tau^*(A) - \tau^*(B)| < \epsilon$, 由于 $\tau^*(B(n)) = 1 \times (p_{10} \times \cdots \times p_{n0}) \neq 0$, 因此 $B = B(n) = (\neg 2q_1) \wedge \cdots \wedge (\neg 2q_n)$ 即为所求.

② 设对 $\forall n \in \mathbb{N}_+$, 都有 $p_{10} \times \cdots \times p_{n0} \geq \epsilon$, 根据引理 1, 数列 $\{p_{n0}\}$ 有一个递增的子数列 $\{p_{k_0}\}$ ($i \in \mathbb{N}_+$, $k \in \mathbb{N}_+$), 则 $\{1 -$

p_{i_0} 为递减数列。根据引理 2, 存在 $n_0 \in \mathbb{N}_+$, 使得数列 $\{1 - p_{i_0}\}$ 的前 n_0 项之积小于 ε , 这时取 $B = \bar{B}(n_0) = q_{i_1} \wedge \cdots \wedge q_{i_{n_0}}$, 只有 $v(q_{i_1}), \dots, v(q_{i_{n_0}})$ 全不取 0 时, $v(\bar{B}(n_0)) = \frac{1}{3}, \frac{2}{3}, 1$, 否则 $v(\bar{B}(n_0)) = 0$ 。所以 $\tau^*(\bar{B}(n_0)) = \frac{1}{3}\mu(\{v \in X | v(\bar{B}(n_0)) = \frac{1}{3}\}) + \frac{2}{3}\mu(\{v \in X | v(\bar{B}(n_0)) = \frac{2}{3}\}) + \mu(\{v \in X | v(\bar{B}(n_0)) = 1\}) \leq \mu(\{v \in X | v(\bar{B}(n_0)) = \frac{1}{3}\}) + \mu(\{v \in X | v(\bar{B}(n_0)) = \frac{2}{3}\}) + \mu(\{v \in X | v(\bar{B}(n_0)) = 1\}) = (1 - p_{i_0}) \times \cdots \times (1 - p_{i_{n_0}})$, 从而有 $|\tau^*(A) - \tau^*(\bar{B}(n_0))| \leq (1 - p_{i_0}) \times \cdots \times (1 - p_{i_{n_0}}) < \varepsilon$ 。

$\tau^*(\bar{B}(n_0)) \geq \frac{1}{3}\mu(\{v \in X | v(\bar{B}(n_0)) = \frac{1}{3}\}) + \frac{1}{3}\mu(\{v \in X | v(\bar{B}(n_0)) = \frac{2}{3}\}) + \frac{1}{3}\mu(\{v \in X | v(\bar{B}(n_0)) = 1\}) = \frac{1}{3}(1 - p_{i_0}) \times \cdots \times (1 - p_{i_{n_0}}) > 0$, 因此 $B = \bar{B}(n_0) = q_{i_1} \wedge \cdots \wedge q_{i_{n_0}}$ 即为所求。

(3) 设 $0 < \tau^*(A) < 1$, 且 A 中有 k 个原子公式, 记 $A = (q_1, q_2, \dots, q_k)$, 这时取自然数 $m \geq k$, 令 $C(n) = (\neg 2q_{m+1}) \wedge \cdots \wedge (\neg 2q_{m+n})$, $B = B(n) = A \vee C(n)$, 由定理 3 得: $\tau^*(B) = \tau^*(A) + \tau^*(C) - \tau^*(A \wedge C) \leq \tau^*(A) + \tau^*(A) - \tau^*(A)\tau^*(C)$, 所以可得 $|\tau^*(A) - \tau^*(B)| \leq (1 - \tau^*(A))\tau^*(C)$, 当 $v(q_{m+1}) = v(q_{m+2}) = \cdots = v(q_{m+n}) = 0$ 时, $v((\neg 2q_{m+1}) \wedge \cdots \wedge (\neg 2q_{m+n})) = 1$, 否则 $v((\neg 2q_{m+1}) \wedge \cdots \wedge (\neg 2q_{m+n})) = 0$ 。所以 $\tau^*(C) = 1 \times (p_{(m+1)0} \times \cdots \times p_{(m+n)0})$, 因此 $|\tau^*(A) - \tau^*(B)| = (1 - \tau^*(A))\tau^*(C) < \tau^*(C) = 1 \times (p_{(m+1)0} \times \cdots \times p_{(m+n)0})$ 。

按照(2)的方法, 可以得到公式 B 满足 $|\tau^*(A) - \tau^*(B)| < \varepsilon, \tau^*(B) \neq \tau^*(A)$ 。

3 公式的相似度和概率逻辑度量空间

定义 6 设 $A, B \in F(S)$, 令 $\xi^*(A, B) = \tau^*((A \rightarrow B) \wedge (B \rightarrow A))$, 则称 $\xi^*(A, B)$ 为 A, B 的概率相似度。显然, $\xi^*(A, B) = \xi^*(B, A)$ 。

定理 5 设 $A, B, C \in F(S)$

- (1) $\xi^*(A, B) = 1$ 当且仅当 $A \approx B$;
- (2) $\xi^*(A, B) + \xi^*(B, C) - 1 \leq \xi^*(A, C)$ 。

证明: (1) 由定理 1 易证;

(2) 首先设 A, B, C 含有 n 个相同的原子公式 q_1, q_2, \dots, q_n , 则

$$\begin{aligned} \xi^*(A, B) &= \tau^*((A \rightarrow B) \wedge (B \rightarrow A)) = \frac{1}{3}\mu([(A \rightarrow B) \wedge (B \rightarrow A)]_{\frac{1}{3}}] + \frac{2}{3}\mu([(A \rightarrow B) \wedge (B \rightarrow A)]_{\frac{2}{3}}]) + \mu([(A \rightarrow B) \wedge (B \rightarrow A)]_1]) \\ &= \mu(\{v \in X | v(A) = v(B)\}) + \frac{1}{3}(\mu(\{v \in X | v(A) = \frac{1}{3}, v(B) = 1\}) + \mu(\{v \in X | v(A) = \frac{1}{3}, v(B) = \frac{2}{3}\}) + \mu(\{v \in X | v(A) = \frac{2}{3}, v(B) = \frac{1}{3}\}) + \mu(\{v \in X | v(A) = \frac{2}{3}, v(B) = 1\}) + \mu(\{v \in X | v(A) = 1, v(B) = \frac{1}{3}\}) + \mu(\{v \in X | v(A) = 1, v(B) = \frac{2}{3}\}) + \mu(\{v \in X | v(A) = 1, v(B) = 1\})); \end{aligned}$$

$$\begin{aligned} \xi^*(B, C) &= \mu(\{v \in X | v(B) = v(C)\}) + \frac{1}{3}(\mu(\{v \in X | v(B) = \frac{1}{3}, v(C) = 1\}) + \mu(\{v \in X | v(B) = \frac{1}{3}, v(C) = \frac{2}{3}\}) + \mu(\{v \in X | v(B) = \frac{2}{3}, v(C) = \frac{1}{3}\}) + \frac{2}{3}(\mu(\{v \in X | v(B) = \frac{2}{3}, v(C) = 1\}) + \mu(\{v \in X | v(B) = 1, v(C) = \frac{2}{3}\})); \end{aligned}$$

$$\begin{aligned} \xi^*(A, C) &= \mu(\{v \in X | v(A) = v(C)\}) + \frac{1}{3}(\mu(\{v \in X | v(A) = \frac{1}{3}, v(C) = 1\}) + \mu(\{v \in X | v(A) = \frac{1}{3}, v(C) = \frac{2}{3}\}) + \mu(\{v \in X | v(A) = \frac{2}{3}, v(C) = \frac{1}{3}\}) + \frac{2}{3}(\mu(\{v \in X | v(A) = \frac{2}{3}, v(C) = 1\}) + \mu(\{v \in X | v(A) = 1, v(C) = \frac{2}{3}\})). \end{aligned}$$

令 $G_1 = \{v \in X | v(A) = v(B)\}, G_2 = \{v \in X | v(B) = v(C)\}, G_3 = \{v \in X | v(A) = v(C)\}$ 。

$$G_1^1 = \{v \in X | v(A) = \frac{1}{3}, v(B) = 1\}, G_2^1 = \{v \in X | v(B) = \frac{1}{3}, v(C) = 1\}, G_3^1 = \{v \in X | v(A) = \frac{1}{3}, v(C) = 1\};$$

$$G_1^2 = \{v \in X | v(A) = 1, v(B) = \frac{1}{3}\}, G_2^2 = \{v \in X | v(B) = 1, v(C) = \frac{1}{3}\}, G_3^2 = \{v \in X | v(A) = 1, v(C) = \frac{1}{3}\};$$

$$G_1^3 = \{v \in X | v(A) = \frac{2}{3}, v(B) = 1\}, G_2^3 = \{v \in X | v(B) = \frac{2}{3}, v(C) = 1\}, G_3^3 = \{v \in X | v(A) = \frac{2}{3}, v(C) = 1\};$$

$$G_1^4 = \{v \in X | v(A) = 1, v(B) = \frac{2}{3}\}, G_2^4 = \{v \in X | v(B) = 1, v(C) = \frac{2}{3}\}, G_3^4 = \{v \in X | v(A) = 1, v(C) = \frac{2}{3}\};$$

$$G_1^5 = \{v \in X | v(A) = \frac{1}{3}, v(B) = \frac{2}{3}\}, G_2^5 = \{v \in X | v(B) = \frac{1}{3}, v(C) = \frac{2}{3}\}, G_3^5 = \{v \in X | v(A) = \frac{1}{3}, v(C) = \frac{2}{3}\};$$

$$G_1^6 = \{v \in X | v(A) = \frac{2}{3}, v(B) = \frac{1}{3}\}, G_2^6 = \{v \in X | v(B) = \frac{2}{3}, v(C) = \frac{1}{3}\}, G_3^6 = \{v \in X | v(A) = \frac{2}{3}, v(C) = \frac{1}{3}\}。$$

从上可以看出 $G_1 \cap G_2 \subset G_3$, 故 $\mu(G_1 \cap G_2) \leq \mu(G_3)$, 又 $\mu(G_1 \cup G_2) = \mu(G_1) + \mu(G_2) - \mu(G_1 \cap G_2)$, 且 $\mu(G_1 \cup G_2) \leq 1$, 即可得 $\mu(G_1) + \mu(G_2) - 1 \leq \mu(G_3)$ 。

又 $G_1 \cap G_2^1 = \emptyset$, 故 $\mu(G_1 \cup G_2^1) = \mu(G_1) + \mu(G_2^1)$ 且 $\mu(G_1 \cup G_2^1) \leq 1$, 从而可得 $\mu(G_1) + \mu(G_2^1) - 1 \leq \mu(G_3^1)$; 同理可得 $\mu(G_2^1) + \mu(G_2^2) - 1 \leq \mu(G_3^2), \mu(G_1^1) + \mu(G_2^2) - 1 \leq \mu(G_3^3), \mu(G_1^1) + \mu(G_2^3) - 1 \leq \mu(G_3^4), \mu(G_1^2) + \mu(G_2^4) - 1 \leq \mu(G_3^5), \mu(G_1^2) + \mu(G_2^5) - 1 \leq \mu(G_3^6)$, 故可得 $\xi^*(A, B) + \xi^*(B, C) - 1 \leq \xi^*(A, C)$ 。

定理 6 设 $A, B \in F(S)$, 令 $\rho^*(A, B) = 1 - \xi^*(A, B)$, 则 $\rho^*(A, B)$ 是 $F(S)$ 上的伪距离, 且 $(F(S), \rho^*)$ 中没孤立点, 称 $(F(S), \rho^*)$ 为概率逻辑度量空间。

证明:由定义 6 和定理 5 可得 $\rho^*(A, B)$ 是 $F(S)$ 上的伪距离。下面证明 $(F(S), \rho^*)$ 中没有孤立点。设 $A \in F(S)$, ε 是任意给定的正数, 以下证明存在公式 $B \in F(S)$, $B \neq A$, 使得 $\rho^*(A, B) < \varepsilon$ 。

设 $A = (q_1, q_2, \dots, q_k)$, 取自然数 $m \geq k$, 令 $C(n) = 2q_{m+1} \vee 2q_{m+2} \vee \dots \vee 2q_{m+n}$, $B = A \wedge C$, 则以下各式成立: $(A \rightarrow B) \wedge (B \rightarrow A) \approx A \rightarrow B \approx A \rightarrow C$; $\xi^*(A, B) = \tau^*(A \rightarrow C) \geq \tau^*(C)$, 因为只有 $v(q_{m+1}) = v(q_{m+2}) = \dots = v(q_{m+n}) = 0$ 时, $v((2q_{m+1}) \vee (2q_{m+2}) \vee \dots \vee (2q_{m+n})) = 0$, 否则 $v((2q_{m+1}) \vee (2q_{m+2}) \vee \dots \vee (2q_{m+n})) = 1$ 。所以 $\tau^*(C) = \tau^*((2q_{m+1}) \vee (2q_{m+2}) \vee \dots \vee (2q_{m+n})) = 1 - 1 \times (p_{(m+1)0} \times \dots \times p_{(m+n)0})$ 。

因此 $\rho^*(A, B) = 1 - \xi^*(A, B) \leq 1 - p_{(m+1)0} \times \dots \times p_{(m+n)0}$ 。

①假设存在 n , 使得 $p_{(m+1)0} \times \dots \times p_{(m+n)0} < \varepsilon$, 则 $B = B(n) = A \wedge C(n) = A \wedge ((2q_{m+1}) \vee \dots \vee (2q_{m+n}))$, 但与 A 不同, 且 $\rho^*(A, B) < \varepsilon$, 故 $B = B(n)$ 为所求。

②假设对 $\forall n \in \mathbb{N}_+$, 都有 $p_{(m+1)0} \times \dots \times p_{(m+n)0} \geq \varepsilon$, 根据引理 1, 数列 $\{p_{(m+i)0}\}$ 有一个递增的子数列 $\{p_{(m+i_k)0}\}$ ($i \in \mathbb{N}_+$, $k \in \mathbb{N}_+$), 则 $\{1 - p_{(m+i_k)0}\}$ 为递减数列。根据引理 2, 存在 $n_0 \in \mathbb{N}_+$, 使得数列 $\{1 - p_{(m+i_k)0}\}$ 的前 n_0 项之积小于 ε , 这时取 $C = \bar{C}(n_0) = (\neg 2q_{(m+i_1)0}) \vee \dots \vee (\neg 2q_{(m+i_{n_0})0})$, 只有 $v(q_{(m+i_1)0}), \dots, v(q_{(m+i_{n_0})0})$ 全不取 0 时 $v(\bar{C}(n_0)) = 0$, 否则 $v(\bar{C}(n_0)) = 1$ 。所以 $\tau^*(\bar{C}(n_0)) = \frac{1}{3} \mu(\{v \in X | v(\bar{C}(n_0)) = \frac{1}{3}\}) + \frac{2}{3} \mu(\{v \in X | v(\bar{C}(n_0)) = \frac{2}{3}\}) + \mu(\{v \in X | v(\bar{C}(n_0)) = 1\}) = \mu(\{v \in X | v(\bar{C}(n_0)) = 1\}) = 1 - (1 - p_{(m+i_1)0}) \dots (1 - p_{(m+i_{n_0})0})$ 。

$\rho^*(A, B) = 1 - \tau^*(A, B) \leq 1 - \tau^*(\bar{C}(n_0)) = (1 - p_{(m+i_1)0}) \times \dots \times (1 - p_{(m+i_{n_0})0}) < \varepsilon$, $\bar{B}(n_0) = A \wedge \bar{C}(n_0) = A \wedge ((\neg 2q_{(m+i_1)0}) \vee \dots \vee (\neg 2q_{(m+i_{n_0})0}))$ 与 A 不同, 故 $B = B(n)$ 为所求。

例 2 设 $A = q_1, B = q_2$, 计算 $\xi^*(A, B), \rho^*(A, B)$ 。

设 $p_{10} = 0.2, p_{1\frac{1}{3}} = 0.3, p_{1\frac{2}{3}} = 0.4, p_{11} = 0.1, p_{20} = 0.3, p_{2\frac{1}{3}} = 0.3, p_{2\frac{2}{3}} = 0.2, p_{21} = 0.2$ 。

解: $\xi^*(A, B) = \tau^*((q_1 \rightarrow q_2) \wedge (q_2 \rightarrow q_1)) = \frac{1}{3} \mu(\{v \in X | v((q_1 \rightarrow q_2) \wedge (q_2 \rightarrow q_1)) = \frac{1}{3}\}) + \frac{2}{3} \mu(\{v \in X | v((q_1 \rightarrow q_2) \wedge (q_2 \rightarrow q_1)) = \frac{2}{3}\}) + \mu(\{v \in X | v((q_1 \rightarrow q_2) \wedge (q_2 \rightarrow q_1)) = 1\}) = \frac{1}{3} (\mu(\{v \in X | v(q_1 \rightarrow q_2) = \frac{1}{3}, v(q_2 \rightarrow q_1) = 1\}) + \mu(\{v \in X | v(q_1 \rightarrow q_2) = 1, v(q_2 \rightarrow q_1) = \frac{1}{3}\}) + \mu(\{v \in X | v(q_1 \rightarrow q_2) = v(q_2 \rightarrow q_1) = \frac{1}{3}\}) + \mu(\{v \in X | v(q_1 \rightarrow q_2) = \frac{2}{3}, v(q_2 \rightarrow q_1) = \frac{2}{3}\}) + \mu(\{v \in X | v(q_1 \rightarrow q_2) = \frac{2}{3}, v(q_2 \rightarrow q_1) = 1\}) + \mu(\{v \in X | v(q_1 \rightarrow q_2) = 1, v(q_2 \rightarrow q_1) = \frac{2}{3}\}) + \mu(\{v \in X |$

$v(q_1 \rightarrow q_2) = v(q_2 \rightarrow q_1) = \frac{2}{3}\}) + \mu(\{v \in X | v(q_1) = v(q_2) = 1\}) = \frac{1}{3} (\mu(\{v \in X | v(q_1) = 1, v(q_2) = \frac{1}{3}\}) + \mu(\{v \in X | v(q_1) = \frac{1}{3}, v(q_2) = 1\}) + \mu(\{v \in X | v(q_1) = \frac{2}{3}, v(q_2) = \frac{2}{3}\}) + \mu(\{v \in X | v(q_1) = \frac{2}{3}, v(q_2) = \frac{1}{3}\}) + \frac{2}{3} (\mu(\{v \in X | v(q_1) = 1, v(q_2) = \frac{2}{3}\}) + \mu(\{v \in X | v(q_1) = \frac{2}{3}, v(q_2) = 1\}) + \mu(\{v \in X | v(q_1) = v(q_2) = 0\}) + \mu(\{v \in X | v(q_1) = v(q_2) = \frac{1}{3}\}) + \mu(\{v \in X | v(q_1) = v(q_2) = \frac{2}{3}\}) + \mu(\{v \in X | v(q_1) = v(q_2) = 1\}) = \frac{1}{3} \times (0.1 \times 0.3 + 0.3 \times 0.2 + 0.3 \times 0.2 + 0.4 \times 0.3) + \frac{2}{3} \times (0.1 \times 0.2 + 0.4 \times 0.2) + 0.2 \times 0.3 + 0.3 \times 0.3 + 0.4 \times 0.2 + 0.1 \times 0.2 = \frac{61}{150} \rho^*(A, B) = 1 - \xi^*(A, B) = 1 - \frac{61}{150} = \frac{89}{150}$ 。

结束语 在四值 Gödel 命题逻辑系统中研究了原子公式

$q_i (i=1, 2, \dots)$ 随机赋值为 $0, \frac{1}{3}, \frac{2}{3}, 1$ 的概率分别为 $p_{i0}, p_{i\frac{1}{3}}, p_{i\frac{2}{3}}, p_{i1}$ ($p_{i0} + p_{i\frac{1}{3}} + p_{i\frac{2}{3}} + p_{i1} = 1$) 的情形, 体现了概率分布的随机性。此外, 利用公式的概率真度定义了公式的概率相似度和伪距离, 建立了概率逻辑度量空间, 为研究该逻辑系统的近似推理奠定了坚实的基础。

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